Abstract—We present a new type of attack in which source code is maliciously encoded so that it appears different to a compiler and to the human eye. This attack exploits subtleties in text-encoding standards such as Unicode to produce source code whose tokens are logically encoded in a different order from the one in which they are displayed, leading to vulnerabilities that cannot be perceived directly by human code reviewers. ‘Trojan Source’ attacks, as we call them, pose an immediate threat both to first-party software and of supply-chain compromise across the industry. We present working examples of Trojan-Source attacks in C, C++, C#, JavaScript, Java, Rust, Go, and Python. We propose definitive compiler-level defenses, and describe other mitigating controls that can be deployed in editors, repositories, and build pipelines while compilers are upgraded to block this attack.

Index Terms—security, compilers, encodings, Unicode

I. INTRODUCTION

What if it were possible to trick compilers into emitting binaries that did not match the logic visible in source code? We demonstrate that this is not only possible for a broad class of modern compilers, but easily exploitable.

We show that subtleties of modern expressive text encodings, such as Unicode, can be used to craft source code that appears visually different to developers and to compilers. The difference can be exploited to invisibly alter the logic in an application and introduce targeted vulnerabilities.

The belief that trustworthy compilers emit binaries that correctly implement the algorithms defined in source code is a foundational assumption of software. It is well-known that malicious compilers can produce binaries containing vulnerabilities [1]; as a result, there has been significant effort devoted to verifying compilers and mitigating their exploitable side-effects. However, to our knowledge, producing vulnerable binaries via unmodified compilers by manipulating the encoding of otherwise non-malicious source code has not so far been explored.

Consider a supply-chain attacker who seeks to inject vulnerabilities into software upstream of the ultimate targets, as happened in the recent Solar Winds incident [2]. Two methods an adversary may use to accomplish such a goal are suborning an insider to commit vulnerable code into software systems, and contributing subtle vulnerabilities into open-source projects. In order to prevent or mitigate such attacks, it is essential for developers to perform at least one code or security review of every submitted contribution. However, this critical control may be bypassed if the vulnerabilities do not appear in the source code displayed to the reviewer, but are hidden in the encoding layer underneath.

Such an attack is quite feasible, as we will hereafter demonstrate.

In this paper, we make the following contributions.

• We define a novel class of vulnerabilities, which we call Trojan-Source attacks, and which use maliciously encoded but semantically permissible source code modifications to introduce invisible software vulnerabilities.
• We provide working examples of Trojan-Source vulnerabilities in C, C++, C#, JavaScript, Java, Rust, Go, and Python.
• We describe effective defenses that must be employed by compilers, as well as other defenses that can be used in editors, repositories, and build pipelines.
• We document the coordinated disclosure process we used to disclose this vulnerability across the industry.
• We raise a new question about what it means for a compiler to be trustworthy.

II. BACKGROUND

A. Compiler Security

Compilers translate high-level programming languages into lower-level representations such as architecture-specific machine instructions or portable bytecode. They seek to implement the formal specifications of their input languages, deviations from which are considered to be bugs.

Since the 1960s [3], researchers have investigated formal methods to mathematically prove that a compiler’s output correctly implements the source code supplied to it [5], [6]. Many of the discrepancies between source code logic and compiler output logic stem from compiler optimizations, about which it can be difficult to reason [7]. These optimizations may also cause side-effects that have security consequences [8].

B. Text Encodings

Digital text is stored as an encoded sequence of numerical values, or code points, that correspond with visual glyphs according to the relevant specification. While single-script
specifications such as ASCII were historically prevalent, modern text encodings have standardized\(^1\) around Unicode \(^9\).

At the time of writing, Unicode defines 143,859 characters across 154 different scripts in addition to various non-script character sets (such as emojis) plus a plethora of control characters. While its specification provides a mapping from numerical code points to characters, the binary representation of those code points is determined by which of various encodings is used, with one of the most common being UTF-8.

Text rendering is performed by interpreting encoded bytes as numerical code points according to the chosen encoding, then looking up the characters in the relevant specification, then resolving all control characters, and finally displaying the glyphs provided for each character in the chosen font.

### C. Supply-Chain Attacks

Supply-chain attacks are those in which an adversary tries to introduce targeted vulnerabilities into deployed applications, operating systems, and software components \(^{10}\). Once published, such vulnerabilities are likely to persist within the affected ecosystem even if patches are later released \(^{11}\). Following a number of attacks that compromised multiple firms and government departments, supply-chain attacks have gained urgent attention from the US White House \(^{12}\).

Adversaries may introduce vulnerabilities in supply-chain attacks through modifying source code, compromising build systems, or attacking the distribution of published software \(^{13}\), \(^{14}\). Distribution attacks are mitigated by software producers signing binaries, so attacks on the earlier stages of the pipeline are particularly attractive. Attacks on upstream software such as widely-utilized packages can affect multiple dependent products, potentially compromising whole ecosystems. As supply-chain threats involve multiple organizations, modeling and mitigating them requires consideration of technical, economic and social factors \(^{15}\).

Open-source software provides a significant vector through which supply-chain attacks can be launched \(^{16}\), and is ranked as one of OWASP’s Top 10 web application security risks \(^{17}\).

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1 According to scans by [w3techs.com/technologies/details/en-utf8](http://w3techs.com/technologies/details/en-utf8), 97% of the most accessed 10 million websites in 2021 use UTF-8 Unicode encodings.

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### III. ATTACK METHODOLOGY

#### A. Reordering

Internationalized text encodings require support for both left-to-right languages such as English and Russian, and right-to-left languages such as Hebrew and Arabic. When mixing scripts with different display orders, there must be a deterministic way to resolve conflicting directionality. For Unicode, this is implemented in the Bidirectional, or Bidi, Algorithm \(^3\).

In some scenarios, the default ordering set by the Bidi Algorithm may not be sufficient; for these cases, override control characters are provided. Bidi overrides are invisible characters that enable switching the display ordering of groups of characters.

Table I provides a list of Bidi override characters relevant to this attack. Of note are LRI and RLI, which format subsequent text as left-to-right and right-to-left respectively, and are both closed by PDI.

Bidi overrides enable even single-script characters to be displayed in an order different from their logical encoding. This fact has previously been exploited to disguise the file extensions of malware disseminated by email \(^{18}\) and to craft adversarial examples for NLP machine-learning pipelines \(^{19}\).

As an example, consider the following Unicode character sequence:

```
RLI a b c PDI
```

which will be displayed as:

```
c b a
```

All Unicode Bidi overrides are restricted to affecting a single paragraph, as a newline character will explicitly close any unbalanced overrides, namely overrides that lack a corresponding closing character.

#### B. Isolate Shuffling

In the Bidi specification, isolates are groups of characters that are treated as a single entity; that is, the entire isolate will be moved as a single block when the display order is overridden.

Isolates can be nested. For example, consider the Unicode character sequence:

```
LRI a b c PDI
```
RLI LRI a b c PDI LRI d e f PDI PDI

which will be displayed as:

d e f a b c

Embedding multiple layers of LRI and RLI within each other enables the near-arbitrary reordering of strings. This gives an adversary fine-grained control, so they can manipulate the display order of text into an anagram of its logically-encoded order.

C. Compiler Manipulation

Like most non-text rendering systems, compilers and interpreters do not typically process formatting control characters, including Bidi overrides, prior to parsing source code. This can be used to engineer a targeted gap between the visually-rendered source code as seen by a human eye, and the raw bytes of the encoded source code as evaluated by a compiler.

We can exploit this gap to create adversarially-encoded text that is understood differently by human reviewers and by compilers.

D. Syntax Adherence

Most well-designed programming languages will not allow arbitrary control characters in source code, as they will be viewed as tokens meant to affect the logic. Thus, randomly placing Bidi override characters in source code will typically result in a compiler or interpreter syntax error. To avoid such errors, we can exploit two general principles of programming languages:

- **Comments** – Most programming languages allow comments within which all text (including control characters) is ignored by compilers and interpreters.
- **Strings** – Most programming languages allow string literals that may contain arbitrary characters, including control characters.

While both comments and strings will have syntax-specific semantics indicating their start and end, these bounds are not respected by Bidi overrides. Therefore, by placing Bidi override characters exclusively within comments and strings, we can smuggle them into source code in a manner that most compilers will accept.

Making a random modification to the display order of characters on a line of valid source code is not particularly interesting, as it is very likely to be noticed by a human reviewer. Our key insight is that we can reorder source code characters in such a way that the resulting display order also represents syntactically valid source code.

E. Novel Supply-Chain Attack

Bringing all this together, we arrive at a novel supply-chain attack on source code. By injecting Unicode Bidi override characters into comments and strings, an adversary can produce syntactically-valid source code in most modern languages for which the display order of characters presents logic that diverges from the real logic. In effect, we anagram program A into program B.

Such an attack could be challenging for a human code reviewer to detect, as the rendered source code looks perfectly acceptable. If the change in logic is subtle enough to go undetected in subsequent testing, an adversary could introduce targeted vulnerabilities without being detected. We provide working examples of this attack in the following section.

Yet more concerning is the fact that Bidi override characters persist through the copy-and-paste functions on most modern browsers, editors, and operating systems. Any developer who copies code from an untrusted source into a protected code base may inadvertently introduce an invisible vulnerability. Such code copying is a significant source of real-world security exploits [20].

F. Generality

We have implemented the above attack methodology, and the examples in the following section, with Unicode. Many modern compilers accept Unicode source code, as will be noted in our experimental evaluation. However, this attack paradigm should work with any text specification that enables the manipulation of display order, which is necessary to support internationalized text. Should the Unicode specification be supplanted by another standard, then in the absence of specific defenses, we believe that it is very likely to provide the same bidirectional functionality used to perform this attack.

IV. EXPLOIT TECHNIQUES

There are a variety of ways to exploit the adversarial encoding of source code. The underlying principle is the same in each: use Bidi overrides to create a syntactically valid reordering of source code characters in the target language.

In the following section, we propose three general types of exploits that work across multiple languages. We do not claim that this list is exhaustive.
A. Early Returns

In the early-return exploit technique, adversaries disguise a genuine return statement as a comment or string literal, so they can cause a function to return earlier than it appears to.

Consider, for example, the case of docstrings – formal comments that purport to document the purpose of a function – which are considered good practice in software development. In languages where docstrings can be located within a function definition, an adversary need only find a plausible location to write the word `return` (or its language-specific equivalent) in a docstring comment, and then reorder the comment such that the `return` statement is executed immediately following the comment.

Figures 1 and 2 depict the encoded bytes and rendered text, respectively, of an early-return attack in Python3. Viewing the rendered text of the source code in fig. 2, one would expect the value of `bank['alice']` to be 50 after program execution. However, the value of `bank['alice']` remains 100 after the program executes. This is because the word `return` in the docstring is actually executed due to a Bidi override, causing the function to return prematurely and the code which subtracts value from a user’s bank account to never run.

This technique is not specific to docstrings; any comment or string literal that can be manipulated by an adversary could hide an early-return statement.

B. Commenting-Out

In this exploit technique, text that appears to be legitimate code actually exists within a comment and is thus never executed. This allows an adversary to show a reviewer some code that appears to be executed but is not present from the perspective of the compiler or interpreter. For example, an adversary can comment out an important conditional, and then reorder the comment such that the `return` statement is executed immediately following the comment.

Figures 3 and 4 depict the encoded bytes and rendered text, respectively, of a commenting-out attack in C. Viewing the rendered text makes it appear that, since the user is not an admin, no text should be printed. However, upon execution the program prints `You are an admin`. The conditional does not actually exist; in the logical encoding, its text is wholly within the comment. This example is aided by the Unicode feature that directionality-aware punctuation characters, such as `, are displayed in reverse within right-to-left settings.

C. Stretched Strings

In this exploit technique, text that appears to be outside a string literal is actually located within it. This allows an adversary to manipulate string comparisons, for example causing strings which appear identical to give rise to a non-equal comparison.

Figures 5 and 6 depict the encoded bytes and rendered text, respectively, of a stretched-string attack in JavaScript. While it appears that the user’s access level is "user" and therefore nothing should be written to the console, the code in fact outputs `You are an admin`. This is because the apparent comment following the comparison isn’t actually a comment, but included in the comparison’s string literal.

In general, the stretched-strings technique will allow an adversary to cause string comparisons to fail.

However, there are other, perhaps simpler, ways that an adversary can cause a string comparison to fail without visual effect. For example, the adversary can place invisible characters – that is, characters in Unicode that render to the absence of a glyph – such as the Zero Width Space (ZWSP) into string literals used in comparisons. Although these invisible characters do not change the way a string literal renders, they will cause string comparisons to fail. Another option is to use characters that look the same, known as homoglyphs, such as the Cyrillic letter ‘х’ which typically renders identical to the Latin letter ‘x’ used in English but occupies a different code point. Depending on the context, the use of other character-encoding tricks may be more desirable than a stretched-string attack using Bidi overrides.

V. RELATED WORK

A. URL Security

Deceptively encoded URLs have long been a tool of choice for spammers [21], with one of the earliest documented examples being the case of `paypal.com`. This July 2000 campaign sought to trick users into disclosing passwords for `paypal.com`
by registering a domain with the lowercase l replaced with the visually similar uppercase I [22].

These domain attacks become even more severe with the introduction of Unicode, which has a much larger set of visually similar characters, or homoglyphs, than ASCII. In fact, Unicode produces a security report which spends considerable length discussing domain-related concerns [23], and the topic of homoglyphs in URLs has been thoroughly examined in the literature [24]–[27].

Punycode, a standard for converting Unicode URLs to ASCII, was created to minimize the attack surface for URL homoglyph attacks [28]. This standard maps well-known homoglyphs to the same Punycode sequences, and prevents the registering of many visually identical URLs.

B. Adversarial NLP

Bidi overrides and homoglyphs have both been used to create adversarial examples in the machine learning NLP setting [19]. These characters, together with invisible characters such as zero-width spaces and deletions control characters, are used to generate strings that look visually identical to some target string but are represented by different Unicode encodings. Optimal encodings are discovered using a gradient-free optimization method that can be used to manipulate the performance of models in both a targeted and untargeted fashion.

C. Visually Deceptive Malware

Bidi overrides have historically been used in the wild to change the appearance of file extensions [18]. This technique aids email-based distribution of malware, as it can deceive a user into running an executable file when they believe they are opening something more benign.

Similarly, directionality overrides have been used in at least one family of malware to disguise the names of malicious system services [29].

Attacks have also been proposed in which an adversary uses homoglyphs to create filenames that look visually similar to key system files, and then replaces references to those files with the adversarial homoglyph version [30].

VI. EVALUATION

A. Experimental Setup

To validate the feasibility of the attacks described in this paper, we have implemented proof-of-concept attacks on simple programs in C, C++, C#, JavaScript, Java, Rust, Go, and Python. Each proof of concept is a program with source code that, when rendered, displays logic indicating that the program should have no output; however, the compiled version of each program outputs the text ‘You are an admin.’ due to Trojan-Source attacks using Bidi override encodings.

For this attack paradigm to work, the compilers or interpreters used must accept some form of Unicode input, such as UTF-8. We find that this is true for the overwhelming majority of languages in modern use. It is also necessary for the language to syntactically support modern internationalized text in string literals or comments.

Future compilers and interpreters should employ defenses that emit errors or warnings when this attack is detected, but we found no evidence of such behavior in any of the experiments we conducted before starting the disclosure process.

All proofs of concept referenced in this paper have been made available online [3]. We have also created a website to help disseminate knowledge of this vulnerability pattern to all developer communities [4].

The following sections describe and evaluate Trojan-Source attack proofs-of-concept against specific programming languages.

B. C

In addition to supporting string literals, C supports both single-line and multi-line comments [31]. Single-line comments begin with the sequence // and are terminated by a newline character. Multi-line comments begin with the sequence /* and are terminated with the sequence */. Conveniently, multi-line comments can begin and end on a single line, despite their name. Strings literal are contained within double quotes, e.g. "· ". Strings can be compared using the function strcmp, which returns a falsey value when strings are equal, and a truty value when strings are unequal.

As previously discussed, Figures 3 and 4 depict a commenting-out attack in C. We also provide an example of a Stretched-String attack in C in Appendix E Figures 24 and 25.

C is well-suited for the commenting-out and stretched-string exploit techniques, but only partially suited for early returns.

This is because when the multline comment terminator, i.e. */ /, is reordered using a right-to-left override, it becomes */. This provides a visual clue that something is not right. This can be overcome by writing reversible comment terminators as // * /, but this is less elegant and still leaves other visual clues such as the line-terminating semicolon. We provide an example of a functioning but less elegant early-return attack in C in Appendix E Figures 26 and 27 which, although it looks like it prints ‘Hello World!’, in fact prints nothing.
We have verified these attacks succeed on both GNU’s gcc v7.5.0 (on Ubuntu) and Apple clang v12.0.5 (on MacOS).

C. C++

Since C++ is a linguistic derivative of C, it should be no surprise that the same attack paradigms work against the C++ specification [32]. Similar proof-of-concept programs modified to adhere to C++ preferred syntax can be seen in Appendix A Figures 8 to 11.

We have verified that both attacks succeed on GNU’s g++ v7.5.0 (on Ubuntu) and Apple clang++ v12.0.5 (on MacOS).

D. C#

C# is an object-oriented language created by Microsoft that typically runs atop .NET, a cross-platform managed runtime, and is used heavily in corporate settings [33]. C# is vulnerable to the same attack paradigms as C and C++, and we present the same proof-of-concept attacks using C# syntax in Appendix B Figures 12 to 15.

We have verified that both attacks succeed on .NET 5.0 using the dotnet-script interpreter on MacOS.

E. JavaScript

JavaScript, also known as ECMAScript, is an interpreted language that provides in-browser client-side scripting for web pages, and is increasingly also used for server-side web application and API implementations [34]. JavaScript is vulnerable to the same attack paradigms as C, C++, and C#, and we present the same proof-of-concept attacks using JavaScript syntax in Appendix C Figures 20 to 23 and 24 as well as the previously discussed Figures 3 and 5.

We have verified that these attacks work against Node.js v16.4.1 (MacOS), which is a local JavaScript runtime built atop Chrome’s V8 JavaScript Engine.

F. Java

Java is a bytecode-compiled multipurpose language maintained by Oracle [35]. It too is vulnerable to the same attack paradigms as C, C++, C#, and JavaScript, and we present the same proof-of-concept attacks using Java syntax in Appendix D Figures 28 to 30.

We have verified that these attacks work against OpenJDK v16.0.1 on MacOS.

G. Rust

Rust is a high-performance language increasingly used in systems programming [36]. It too is vulnerable to the same attack paradigms as C, C++, C#, JavaScript, and Java, and we present the same proof-of-concept attacks using Rust syntax in Appendix E Figures 32 to 34.

We have verified that these attacks work against Rust v1.53.0 (on MacOS), but note that one of the two proof-of-concept (depicted in Figures 22 and 23) throws an unused variable warning on compilation. However, this warning could be evaded by an adversary using the warned variable trivially elsewhere in the program.

H. Go

Go is a multipurpose open-source language produced by Google [37]. Go is vulnerable to the same attack paradigms as C, C++, C#, JavaScript, Java, and Rust, and we present the same proof-of-concept attacks using Go syntax in Appendix F Figures 28 and 29.

We have verified that these attacks work against Go v1.16.6 on MacOS. We note that unused variables throw compiler errors in the official Go compiler, and so our commenting-out Trojan-Source attack proof-of-concept deviates from our general pattern to ensure that no variables are left unused.

I. Python

Python is a general-purpose scripting language used heavily in data science and many other settings [38]. Python supports multiline comments in the form of docstrings opened and closed with ‘’’ or ””. We have already exploited this fact in Figures 1 and 2 to craft elegant early-return attacks.

An additional commenting-out proof-of-concept attack against Python 3 can be found in encoded form in Appendix H Figures 34 and 35.

We have verified that these attacks work against Python 3.9.5 compiled using clang 12.0.0 (on MacOS) and against Python 3.7.10) compiled using GNU’s gcc (on Ubuntu).

J. Code Viewers

We were curious to see how these attacks were visualized by the editors and code repository front-ends used in modern development environments, as many tools have different Unicode implementations. We therefore tested the latest releases of the Visual Studio Code, Atom, Sublime Text, Notepad++, vim, and emacs code editors. We also tested the GitHub and Bitbucket web-based code repository front-end interfaces. Each evaluation was repeated across three machines running Windows 10, MacOS Big Sur, and Ubuntu 20.04. The results can be found in Appendix Table II, where ✓ represents code that displayed the same as the example visualizations in this paper as of the time of writing. Any deviations from this display are described.

VII. DISCUSSION

A. Ethics

We have followed our department’s ethical guidelines carefully throughout this research. We did not launch any attacks using Trojan-Source methods against codebases we did not own. Furthermore, we made responsible disclosure to all companies and organizations owning products in which we discovered vulnerabilities. We offered a 99-day embargo period following our first disclosure to allow affected products to be repaired, which we will discuss later.

B. Attack Feasibility

Attacks on source code are both extremely attractive and highly valuable to motivated adversaries, as maliciously inserted backdoors can be incorporated into signed code that persists in the wild for long periods of time. Moreover, if
backdoors are inserted into open source software components that are included downstream by many other applications, the blast radius of such an attack can be very large. Trojan-Source attacks introduce the possibility of inserting such vulnerabilities into source code invisibly, thus completely circumventing the current principal control against them, namely human source code review. This can make backdoors harder to detect and their insertion easier for adversaries to perform.

There is a long history of the attempted insertion of backdoors into critical code bases. One example was the attempted insertion of a root user escalation-of-privilege backdoor into the Unix kernel, which was as subtle as changing an == token to an = token [39]. This attack was detected when experienced developers saw the vulnerability. The techniques described here allow a similar attack in the future to be invisible.

Recent research in developer security usability has documented that a significant portion of developers will gladly copy and paste insecure source code from unofficial online sources such as Stack Overflow [20], [40]. Since Bidi overrides persist through standard copy-and-paste functionality, malicious code snippets with invisible vulnerabilities can be posted online in the hope that they will end up in production code. The market for such vulnerabilities is vibrant, with exploits on major platforms now commanding seven-figure sums [41].

Our experiments indicate that, as of the time of writing, C, C++, C#, JavaScript, Java, Rust, Go, and Python are all vulnerable to Trojan-Source attacks. More broadly, this class of attacks is likely applicable to any language with common compilers that accept Unicode source code. Any entity whose security relies on the integrity of software supply chains should be concerned.

C. Syntax Highlighting

Many developers use text editors that, in addition to basic text editing features, provide syntax highlighting for the languages in which they are programming. Moreover, many code repository platforms, such as GitHub [2], provide syntax highlighting through a web browser. Comments are often displayed in a different color from code, and many of the proofs of concept provided in this paper work by deceiving developers into thinking that comments are code or vice versa.

We might have hoped that a well-implemented syntax-highlighting platform would at the very least exhibit unusual syntax highlighting in the vicinity of Bidi overrides in code, but our experience was mixed. Some attacks provided strange highlighting in a subset of editors, which may suffice to alert developers that an encoding issue is present. However, all syntax highlighting nuances were editor-specific, and other attacks did not show abnormal highlighting in the same settings.

Although unexpected coloring of source code may flag the possibility of an encoding attack to experienced developers, and in particular to those familiar with this work, we expect that most developers would not even notice unusual highlighting, let alone investigate it thoroughly enough to work out what was going on. A motivated attacker could experiment with the visualization of different attacks in the text editors and code repository front-ends used in targeted organizations in order to select an attack that has no or minimal visual effect.

Bidi overrides will typically cause a cursor to jump positions on a line when using arrow keys to click through tokens, or to highlight a line of text character-by-character. This is an artifact of the effect of the logical ordering of tokens on many operating systems and Unicode implementations. Such behavior, while producing no visible changes in text, may also be enough to alert some experienced developers. However, we suspect that this requires more attention than is given by most developers to reviews of large pieces of code.

D. Invisible Character Attacks

When discussing the string-stretching technique, we proposed that invisible characters or homoglyphs could be used to make visually-identical strings that are logically different when compared. Another invisible-vulnerability technique with with experimentations – largely without success – was the use of invisible characters in function names.

We theorized that invisible characters included in a function name could define a different function from the function defined by only the visible characters. This could allow an attacker to define an adversarial version of a standard function, such as printf in C, that can be invoked by calling the function with an invisible character in the function name. Such an adversarial function definition could be discreetly added to a codebase by defining it in a common open-source package that is imported into the global namespace of the target program.

However, we found that all compilers analyzed in this paper emitted compilation errors when this technique was employed, with the exception of one compiler – Apple clang v12.0.5 – which emitted a warning instead of an error.

Should a compiler not instrument defenses against invisible characters in function definition names – or indeed in variable names – this attack may well be feasible. That said, our experimental evidence suggests that this theoretical attack already has defenses employed against it by most modern compilers, and thus is unlikely to work in practice.

E. Homoglyph Attacks

After we investigated invisible characters, we wondered whether homoglyphs in function names could be used to define distinct functions whose names appeared to the human eye to be the same. Then an adversary could write a function whose name appears the same as a pre-existing function – except that one letter is replaced with a visually similar character. Indeed, this same technique could be used on code identifiers of any kind, such as variables and class names, and may be particularly insidious for homoglyphs that appear like numbers.
We were able to successfully implement homoglyph attack proofs-of-concept in every language discussed in this paper; that is, C, C++, C#, JavaScript, Java, Rust, Go, and Python all appear to be vulnerable. In our experiments, we defined two functions that appeared to have the name sayHello, except that one version used a Latin H while the other used a Cyrillic H.

Consider Figure 7 which implements a homoglyph attack in C++. For clarity, we denote the Latin H in blue and the Cyrillic H in red. This program outputs the text Goodbye, World! when compiled using clang++. Although this example program appears harmless, a homoglyph attack could cause significant damage when applied against a common function, perhaps via an imported library. For example, suppose a function called hashPassword was replaced with a similar function that called and returned the same value as the original function, but only after leaking the pre-hashed password over the network.

All compilers and interpreters examined in this paper emitted the text Goodbye, World! with similar proofs of concept. There were only three exceptions. GNU’s gcc and its C++ counterpart, g++, both emitted stray tokens errors. Of particular note is the Rust compiler, which threw a ‘mixed_script_confusables’ warning while producing the homoglyph attack binary. The warning text suggested that the function name with the Cyrillic Hused “mixed script confusables” and suggested rechecking to ensure usage of the function was wanted. This is a well-designed defense against homoglyph attacks, and it shows that this attack has been seriously considered by at least one compiler team.

This defense, together with the defenses against invisible character attacks, should serve as a precedent. It is reasonable to expect compilers to also incorporate defences against Trojan-Source attacks.

F. Defenses

The simplest defense is to ban the use of text directionality control characters both in language specifications and in compilers implementing these languages.

In most settings, this simple solution may well be sufficient. If an application wishes to print text that requires Bidi overrides, developers can generate those characters using escape sequences rather than embedding potentially dangerous characters into source code.

This simple defense can be improved by adding a small amount of nuance. By banning all directionality-control characters, users with legitimate Bidi-override use cases in comments are penalized. Therefore, a better defense might be to ban the use of unterminated Bidi override characters within string literals and comments. By ensuring that each override is terminated – that is, for example, that every LRI has a matching PDI – it becomes impossible to distort legitimate source code outside of string literals and comments.

Trojan-Source defenses must be enabled by default on all compilers that support Unicode input, and turning off the defenses should only be permitted when a dedicated suppression flag is passed.

While changes to language specifications and compilers are ideal solutions, there is an immediate need for existing code bases to be protected against this family of attacks. Moreover, some languages or compilers may choose not to implement appropriate defenses. To protect organizations that rely on them, defenses can be employed in build pipelines, code repositories, and text editors.

Build pipelines, such as those used by software producers to build and sign production code, can scan for the presence of Bidi overrides before initiating each build and break the build if such a character is found in source code. Alternatively, build pipelines can scan for the more nuanced set of unterminated Bidi overrides. Such tactics provide an immediate and robust defense for existing software maintainers.

Code repository systems and text editors can also help prevent Trojan-Source attacks by making them visible to human reviewers. For example, code repository front-ends, such as web UIs for viewing committed code, can choose to represent Bidi overrides as visible tokens, thus making attacks visible, and by adding a visual warning to the affected lines of code.

Code editors can employ similar tactics. In fact, some already do; for example, defaults to showing Bidi overrides as numerical code points rather than applying the Bidi algorithm. However, many common code editors do not adopt this behavior, including most GUI editors such as, at the time of writing, Microsoft’s VS Code and Apple’s Xcode.

G. Coordinated Disclosure

We contacted nineteen independent companies and organizations in a coordinated disclosure effort to build defenses for affected compilers, interpreters, code editors, and code repository front-ends. We set a 99-day embargoed disclosure period during which disclosure recipients could implement defenses before we published our attacks. We met a variety of responses ranging from patching commitments and bug bounties to quick dismissal and references to legal policies.

We selected an initial set of disclosure recipients by identifying the maintainers of products that our experiments indicated were affected by the Trojan Source vulnerability.
pattern. We also included companies that, to our knowledge, maintained their own internal compilers and build tools. The initial disclosures were sent on July 25, 2021.

Several of the initial recipients asked us to include additional organizations in the disclosure process, and we did so. We also sent additional disclosures throughout the embargo window for affected products that we discovered during the disclosure process.

Of the nineteen software suppliers with whom we engaged, seven used an outsourced platform for receiving vulnerability disclosures, six had dedicated web portals for vulnerability disclosures, four accepted disclosures via PGP-encrypted email, and two accepted disclosures only via non-PGP email. They all confirmed receipt of our disclosure, and ultimately nine of them committed to releasing a patch.

Eleven of the recipients had bug bounty programs offering payment for vulnerability disclosures. Of these, five paid bounties, with an average payment of $2,246.40 and a range of $4,475.

On September 9, 2021, we sent a vulnerability report to CERT/CC, the CERT Coordination Center sponsored by CISA [42]. Our report was accepted the same day for coordinated disclosure assistance. This gave all affected vendors access to VINCE, a tool providing a shared communication platform across vendors implementing defenses. Thirteen of our recipients, inclusive of CERT/CC, opted in to the VINCE tool for these shared communications. CERT/CC also added three additional vendors to the disclosure beyond the nineteen we had already contacted.

On October 18, 2021, Trojan Source attacks were issued two CVEs [43]: CVE-2021-42574 for tracking the Bidi attack, and CVE-2021-42694 for tracking the homoglyph attack. These CVEs were issued by MITRE against the Unicode specification.

On the same day, we sent a PGP-encrypted disclosure to the distros mailing list [44], which contains representatives of the security teams of 21 operating systems as of the time of writing. This list coordinates the application of patches across OS maintainers, but allows a maximum embargo period of 14 days.

We observed multiple patterns throughout the coordinated disclosure process:

1) Novel Vulnerability Patterns: Vulnerability disclosures which do not follow commonly known vulnerability patterns (such as CWEs [45]) are likely to be screened out by disclosure recipients. We observed a tendency to close issues immediately as representing no threat when they did not align to something well-known and easily evidenced, such as SQL injection. This was particularly the case with software suppliers that outsource the bug-bounty award process to third-party vulnerability submission services. We found that we could generally restart the disclosure process when this happened by reaching out to personal contacts affiliated with the recipient firm, or failing that by asking the vulnerability submission service to put us in contact with a full-time security team.

2) Impactful Language: When writing vulnerability disclosures, descriptions that personalise the potential impact can be needed to drive action. Neutral disclosures like those found in academic papers are less likely to evoke a response than disclosures stating that named products are immediately at risk.

3) CVEs: CVEs are really useful, as they increase the chance that the recipient will take the time to actually read and understand the report. However, CVEs are by default raised by the affected supplier, so are not much help with the initial contact. We eventually had to fall back on the CVE issuer of last resort, MITRE.

4) Shared Communication: CERT/CC’s VINCE platform provides a useful and neutral cross-organization discussion tool during coordinated disclosures. The tool allows affected vendors to post on a private discussion board, and makes it much easier to communicate to all affected parties in a central location. The CERT/CC team will also help to coordinate contacting affected vendors under embargo, which provides a helpful method for scaling out disclosure efforts at no cost. Like CVEs, having a CERT/CC case also adds to the credibility of disclosures.

5) Open Source Assistance: Disclosing to open source operating system security teams is helpful for assistance coordinating patches across the ecosystem, including with contributors of open source projects that may not otherwise offer an embargoed disclosure method. In particular, Linux operating systems backed by a commercial entity have both the funding and incentives to ensure that common open source tools are patched prior to public disclosure. Maintainers of open source projects commonly work for or closely with these companies, and as such can be included in security responses.

H. Ecosystem Scanning

We were curious if we could find any examples of Trojan Source attacks in the wild prior to public disclosure of the attack vector, and therefore tried to scan as much of the open source ecosystem as we could for signs of attack.

We assembled a regex that identified unterminated Bidi override sequences in comments and strings, and GitHub provided us with the results of this pattern run against all public commits containing non-markup language source code ingested into GitHub from January through mid October 2021. This yielded 7,444 commits, which resolved to 2,096 unique files still present in public repositories as of October 2021.

The majority of the results were false positives. Examples of clearly non-malicious encodings included LRE characters placed at the start of file paths, malformed strings in genuinely right-to-left languages, and Bidi characters placed into localized format string patterns.

However, we did find some evidence of techniques similar to Trojan Source attacks being exploited. In one instance, a static code analysis tool for smart contracts, Slither [46], contained scanning for right-to-left override characters. The tool provides an example of why this scan is necessary: it uses an RLO character to swap the display order of two single-character
variables passed as arguments. In another instance, we discovered the use of RLI and LRI characters used to conceal an invocation of system("cat /etc/passwd"); within a Ruby script. However, the technique used left visual artifacts and appears to be a test of UTF8 functionality, rather than an exploit payload. We also discovered multiple instances of JavaScript obfuscation that used Bidi characters to assist in obscuring code. This is not necessarily malicious, but is still an interesting use of directionality overrides. Finally, we found multiple implementations of exploit generators for directionality override in filename extensions, as previously referenced [18].

In parallel, contributors to the Rust project scanned all historical submissions to crates.io, Rust’s package manager, and found no evidence of exploitation within the Rust ecosystem.

VIII. Conclusion

We have presented a new type of attack that enables invisible vulnerabilities to be inserted into source code. Our Trojan-Source attacks use Unicode control characters to modify the order in which blocks of characters are displayed, thus enabling comments and strings to appear to be code and vice versa. This enables an attacker to craft code that is interpreted one way by compilers and a different way by human reviewers. We present proofs of concept for C, C++, C#, JavaScript, Java, Rust, Go, and Python, and argue that this attack may well appear in any programming language that supports internationalized text in comments and string literals, even in other encoding standards.

As powerful supply-chain attacks can be launched easily using these techniques, it is essential for organizations that participate in a software supply chain to implement defenses. We have discussed countermeasures that can be used at a variety of levels in the software development toolchain: the language specification, the compiler, the text editor, the code repository, and the build pipeline. We are of the view that the long-term solution to the problem will be deployed in compilers. We note that almost all compilers already defend against one related attack, which involves creating adversarial function names using zero-width space characters, while three generate errors in response to another, which exploits homoglyphs in function names.

About half of the compiler maintainers we contacted during the disclosure period are working on patches or have committed to do so. As the others are dragging their feet, it is prudent to deploy other controls in the meantime where this is quick and cheap, or relevant and needful. Three firms that maintain code repositories are also deploying defenses. We recommend that governments and firms that rely on critical software should identify their suppliers’ posture, exert pressure on them to implement adequate defenses, and ensure that any gaps are covered by controls elsewhere in their toolchain.

The fact that the Trojan Source vulnerability affects almost all computer languages makes it a rare opportunity for a system-wide and ecologically valid cross-platform and cross-vendor comparison of responses. As far as we are aware, it is an unprecedented test of the coordinated disclosure ecosystem. However, since the work is still in progress, a full discussion of what we are learning and what might be done better is for a later paper.

Scientifically, this research also contributes to the growing body of work on security usability from the developer’s perspective. It is not sufficient for a compiler to be verified; it must also be safely usable. Compilers that are trivially vulnerable to adversarial text encoding cannot reasonably be described as secure.

Acknowledgment

We would like to thank GitHub for assisting with scanning the open source ecosystem, Pietro Albini and Mara Bos of the Rust project for scanning crates.io, and CERT/CC for assistance with coordinated disclosure.

References


Appendix

A. C++ Trojan-Source Proofs-of-Concept

```cpp
#include <iostream>
#include <string>

int main() {
    std::string access_level = "user";
    if (access_level.compare("user\0 LRI") == 0) {
        std::cout << "You are an admin.\n";
        return 0;
    }
}
```

Fig. 8. Encoded bytes of a Trojan-Source stretched-string attack in C++.

```cpp
#include <iostream>
#include <string>

int main() {
    std::string access_level = "user";
    if (access_level.compare("user") == 0) { // Check if admin
        std::cout << "You are an admin.\n";
        return 0;
    }
}
```

Fig. 9. Rendered text of a Trojan-Source stretched-string attack in C++.

```cpp
#include <iostream>

int main() {
    bool isAdmin = false;
    /* RLO */ LRI
    if (isAdmin) begin admins only */
    std::cout << "You are an admin.\n";
    /* end admin only RLO */
    return 0;
}
```

Fig. 10. Encoded bytes of a Trojan-Source commenting-out attack in C++.

```cpp
#include <iostream>

int main() {
    bool isAdmin = false;
    /* begin admins only */ if (isAdmin) {
        std::cout << "You are an admin.\n";
        /* end admins only */
    }
    return 0;
}
```

Fig. 11. Rendered text of a Trojan-Source commenting-out attack in C++.

B. C# Trojan-Source Proofs-of-Concept

```cs
@/usr/bin/env dotnet-script

string access_level = "user";
if (access_level == "user\0 LRI") {
    Console.WriteLine("You are an admin.");
}
```

Fig. 12. Encoded bytes of a Trojan-Source stretched-string attack in C#.

```cs
@/usr/bin/env dotnet-script

string access_level = "user";
if (access_level == "user") { // Check if admin
    Console.WriteLine("You are an admin.");
}
```

Fig. 13. Rendered text of a Trojan-Source stretched-string attack in C#.

```cs
@/usr/bin/env dotnet-script

bool isAdmin = false;
/* RLO */ LRI
if (isAdmin) begin admins only */
Console.WriteLine("You are an admin.");
/* end admin only RLO */
```

Fig. 14. Encoded bytes of a Trojan-Source commenting-out attack in C#.

```cs
@/usr/bin/env dotnet-script

bool isAdmin = false;
/* begin admins only */ if (isAdmin) {
    Console.WriteLine("You are an admin.");
    /* end admins only */
}
```

Fig. 15. Rendered text of a Trojan-Source commenting-out attack in C#.
C. Java Trojan-Source Proofs-of-Concept

```java
public class TrojanSource {
    public static void main(String[] args) {
        String accessLevel = "user";
        if (accessLevel != "userLRO LRI/\ Check if adminPDI LRI") {
            System.out.println("You are an admin.");
        } /* end admin only RLO { LRI*/
    }
}
```

Fig. 16. Encoded bytes of a Trojan-Source stretched-string attack in Java.

```java
public class TrojanSource {
    public static void main(String[] args) {
        String accessLevel = "user";
        if (accessLevel != "user") { // Check if admin
            System.out.println("You are an admin.");
        }
    }
}
```

Fig. 17. Rendered text of a Trojan-Source stretched-string attack in Java.

```java
public class TrojanSource {
    public static void main(String[] args) {
        boolean isAdmin = false;
        /*RLO */ LRI if (isAdmin) begin admins only */
        System.out.println("You are an admin.");
    } /* end admin only RLO { LRI*/
}
```

Fig. 18. Encoded bytes of a Trojan-Source commenting-out attack in Java.

```java
public class TrojanSource {
    public static void main(String[] args) {
        boolean isAdmin = false;
        /* begin admins only */ if (isAdmin) {
            System.out.println("You are an admin.");
        } /* end admins only */
    }
}
```

Fig. 19. Rendered text of a Trojan-Source commenting-out attack in Java.

D. Rust Trojan-Source Proofs-of-Concept

```rust
fn main() {
    let access_level = "user";
    if (access_level != "userLRO LRI// Check if adminPDI LRI") {
        println!("You are an admin.");
    }
}
```

Fig. 20. Encoded bytes of a Trojan-Source stretched-string attack in Rust.

```rust
fn main() {
    let access_level = "user";
    if access_level != "user" { // Check if admin
        println!("You are an admin.");
    }
}
```

Fig. 21. Rendered text of a Trojan-Source stretched-string attack in Rust.

```rust
fn main() {
    let is_admin = false;
    /*RLO */ LRI if is_adminPDI LRI begin admins only */
    println!("You are an admin.");
    /* end admin only RLO { LRI*/
}
```

Fig. 22. Encoded bytes of a Trojan-Source commenting-out attack in Rust.

```rust
fn main() {
    let is_admin = false;
    /* begin admins only */ if is_admin {
        println!("You are an admin.");
    } /* end admins only */
}
```

Fig. 23. Rendered text of a Trojan-Source commenting-out attack in Rust.
E. C Trojan-Source Proofs-of-Concept

```c
#include <stdio.h>
#include <string.h>

int main() {
    char* access_level = "user";
    if (strcmp(access_level, "user\R\L\I\L\I// Check if admin\P\D\I\ L\I\I") { // Check if admin
        printf("You are an admin\n");
    }
    return 0;
}
```

Fig. 24. Encoded bytes of a Trojan-Source stretched-string attack in C.

```c
#include <stdio.h>

int main() {
    /* Say hello; newline \R\L\I */ return 0;
    printf("Hello world.\n");
    return 0;
}
```

Fig. 25. Rendered text of a Trojan-Source stretched-string attack in C.

```c
#include <stdio.h>
#include <string.h>

int main() {
    char* access_level = "user";
    if (strcmp(access_level, "user") { // Check if admin
        printf("You are an admin\n");
    }
    return 0;
}
```

Fig. 26. Encoded bytes of a Trojan-Source early-return attack in C.

```c
#include <stdio.h>

int main() {
    /* Say hello; newline; return 0 */
    printf("Hello world.\n");
    return 0;
}
```

Fig. 27. Rendered text of a Trojan-Source early-return attack in C.

F. Go Trojan-Source Proofs-of-Concept

```go
package main
import "fmt"

func main() {
    var accessLevel = "user"
    if accessLevel != "user\R\L\I\L\I// Check if admin\P\D\I\ L\I\I" { // Check if admin
        fmt.Println("You are an admin."
    }
}
```

Fig. 28. Encoded bytes of a Trojan-Source stretched-string attack in Go.

```go
package main
import "fmt"

func main() {
    var accessLevel = "user"
    if accessLevel != "user" { // Check if admin
        fmt.Println("You are an admin."
    }
}
```

Fig. 29. Rendered text of a Trojan-Source stretched-string attack in Go.

```go
package main
import "fmt"

func main() {
    var isAdmin = false
    var isSuperAdmin = false
    isAdmin = isAdmin || isSuperAdmin
    /*RL0 \L\I\I if (isAdmin)PD\I L\I\I begin admins only */
    fmt.Println("You are an admin."
    /* end admin only RL0 { L\I\I*/
}
```

Fig. 30. Encoded bytes of a Trojan-Source commenting-out attack in Go.

```go
package main
import "fmt"

func main() {
    var isAdmin = false
    var isSuperAdmin = false
    isAdmin = isAdmin || isSuperAdmin
    /* begin admins only */ if (isAdmin) {
        fmt.Println("You are an admin."
    /* end admins only */
}
```

Fig. 31. Rendered text of a Trojan-Source commenting-out attack in Go.
**G. JavaScript Trojan-Source Proof-of-Concept**

```javascript
#!/usr/bin/env node

var isAdmin = false;
/*@RLO */
LRI if (isAdmin) /* begin admins only */
LRI console.log("You are an admin.");
LRI /* end admin only */ RLO { LRI*/

Fig. 32. Encoded bytes of a Trojan-Source commenting-out attack in JS.
```

```javascript
#!/usr/bin/env node

var isAdmin = false;
/*@RLO */
LRI if (isAdmin) { /* begin admins only */
LRI console.log("You are an admin.");
LRI /* end admins only */ }

Fig. 33. Rendered text of a Trojan-Source commenting-out attack in JS.
```

**H. Python Trojan-Source Proof-of-Concept**

```python
#!/usr/bin/env python3

access_level = "user"
if access_level != "none" or access_level != "user"
    print("You are an admin.");

Fig. 34. Encoded bytes of a Trojan-Source commenting-out attack in Python.
```

```python
#!/usr/bin/env python3

access_level = "user"
if access_level != "none" and access_level != "user": # Check if admin
    print("You are an admin.")

Fig. 35. Rendered text of a Trojan-Source commenting-out attack in Python.
```

**I. Code Viewer Visualizations**

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